Reverse Engineering

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What is Reverse Engineering? Goals of the course

- Rapid immersion: real systems, real tools, real specs, lots of surprising complications
 - You'll need to explore, search (a lot), and prioritize, not just follow the examples
- RE requires some creativity and intuition that is only developed through practice, not just listening to lectures
- You won't be taught everything you need to know in lectures. To be good at RE you have to learn to find answers on your own. Google is your friend.
- By the end of the course you should be able to take a blob of code and figure out what it does. (within reason)
- You will be exposed to malicious software (malware) so you will gain an understanding of what you have seen in the news about ransomware, etc

What is Reverse Engineering? Homework

- There will be homework after every class that is due prior to the next one
 - Tuesday homework will be easier and are designed to solidify what you learned that day
 - Thursday homework will be a little more complex
- Many small assignments are better than fewer large ones. Getting better at RE is like learning a language. A little every day is better than a lot crammed just before it is due.

What is Reverse Engineering? Mid-term & Final

- The course is project-based & all exams will be take-home
 - Subject to the Dartmouth Honor Code
- You will be given some software that you will RE and write up what it does
 - Probably some other small tasks
- If you want to go above and beyond the coursework, individual research projects and Senior Honors Theses will be encouraged
 - Improving state-of-the-art RE tools such as Ghidra and Binary Ninja is strongly encouraged

What is Reverse Engineering? And why do we do it?

- Engineers work from a Design to an Artifact
- Reverse Engineers work backward, from an Artifact back to a Design
- In computing, this is often working from an Executable back to Source Code
- This is useful for many reasons:
 - We can preserve and sustain old software by emulating it
 - We can find security bugs without source code
 - We can copy software, or determine whether one program copies another

What is Reverse Engineering? And is it legal?

- This isn't a course in law, nor are any of us law experts. Seek your own legal advice.
- There are many legal uses for Reverse Engineering, but also there are potential violations of law or contracts.
- The Electronic Frontier Foundation (EFF) has a helpful guide for reference at https://www.eff.org/issues/coders/ reverse-engineering-faq
- "Five areas of United States law are particularly relevant for computer scientists engaging in reverse engineering: • Copyright law and fair use, codified at 17 U.S.C. 107;
- - Trade secret law;
 - The anti-circumvention provisions of the Digital Millennium Copyright Act (DMCA), codified at 17 U.S.C. section 1201;
 - Contract law, if use of the software is subject to an End User License Agreement (EULA), Terms of Service notice (TOS), Terms of Use notice (TOU), Non-Disclosure Agreement (NDA), developer agreement or API agreement; and
 - The Electronic Communications Privacy Act, codified at 18 U.S.C. 2510 et. seq." (-EFF)

We didn't talk about this in class, but we will.



What is Reverse Engineering? And how is it done?

- Source Code is compiled and linked into Machine Code.
- Machine Code looks like this:
 - E8 F9 CA AD DE
- Machine Code translates directly to Assembly Code, like this:
 - CALL 0xDEADCAFE
- At a low level, we're just reading disassembly and annotating it to be legible
- At a high level, we're also trying to understand the program design



What is Reverse Engineering? And how is it done?

- You can learn to read Disassembly, but there are complications:
 - It is very verbose, much more so than C
 - It often lacks variable and function names
- Tools can help!
 - Decompiling the Disassembly into C, or something like C
 - Accepting new variable and function names
 - Transferring symbol names between different programs

What is Reverse Engineering? And what tools make it easier?

- IDA Pro
 - First popular Interactive Disassembler
 - First commercially useful Decompiler
- Ghidra
- **Binary Ninja**
 - Commercial disassembler with clean scripting

We will use Ghidra extensively.

Your final project or Honors Thesis could be a Ghidra plugin or another **Ghidra** improvement

NSA's internal tool for reverse engineering, now with a declassified & free version

IDA - Z:\bir	n\bash
<u>File E</u> dit <u>J</u> ump Searc <u>h V</u> iew Deb <u>ugg</u> er <u>O</u> ptions <u>W</u> indows Help	
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f sub_100002C95	movzx eax, byte ptr [rdi]
text:00000000000000000000000000000000000	cmp eax. 9
<i>f</i> sub_100002D70	jz s 🕘 🔵 About
	test a
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sub_100002F5D	jz s ball
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text:0000000100002D0E	
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sub_1000031BA	iz s
sub_10000320C	movzx e
sub_100003287	cmp e
	jnz s
text:00000336A	www.hex-rays.com
sub_1000033BF	
Text:000000000000000000000000000000000000	
Line 16 of 1768	nchronized w OK <u>A</u> ddons
Output window	□æ×
The initial autoanalysis has been finished.	
Python	
AU: idle Down Disk: 424GB	

The Basics of Reverse Engineering What it's like inside a C program.

air% cat pointers.c #include <stdio_h> #include <stdint_h> #include <stdlib_h>

int main(int argc, char **argv){ void *heapthing=malloc(512); printf("Main is at 0x%08llx.\n", (uint64_t) main); printf("The call stack is near 0x%08llx.\n", (uint64_t) &argc); printf("The heap is near 0x%08llx.\n", (uint64 t) heapthing);

air% ./pointers Main is at 0x100003e84. The call stack is near 0x16fdff76c. The heap is near 0x100304310. air%

Run this program several times. If you see different addresses between runs, you are seeing the effects of ASLR.

Note which part of each address doesn't change: that's because ASLR is at page granularity (typically 4Kbytes)



The Basics of Reverse Engineering

What does Disassembly look like?

Note Intel syntax of this disassembler

00101	169 -	main	2	🖌 - 🍓 🗔 🕴
undefin	ed ma	in()		
[undef	ined	AL:1	<return></return>
	undef	ined8	Stack[-0x	10]:8 local_10
	undef	ined4	Stack[-0x	<pre>1c]:4 local_1c</pre>
	undef	ined8	Stack[-0x	28]:8 local_28
		main		
1169	END			
116d	PUSH	RBP		
116e	MOV	RBP, RSP		
1171	SUB	RSP,0x20		
1175	MOV	dword ptr [R	<pre>BP + local</pre>	_1c]
1178	MOV	qword ptr [R	<pre>BP + local</pre>	_28]
117c	MOV	EDI,0x200		
1181	CALL	<external>::r</external>	nalloc	
1186	MOV	qword ptr [R	<pre>BP + local</pre>	_10]
118a	LEA	RAX,[main]		
1191	MOV	RSI=>main,RA	< Contract of the second secon	
1194	LEA	RDI,[s_Main_:	is_at_0x%0	8llx
119b	MOV	EAX,0x0		
11a0	CALL	<external>::</external>	printf	
11a5	LEA	RAX=>local_1	c,[RBP + -	0x14]
11a9	MOV	RSI,RAX		
11ac	LEA	RDI,[s_The_ca	all_stack_	is_n
11b3	MOV	EAX,0x0		
11b8	CALL	<external>::;</external>	printf	
11bd	MOV	RAX, qword pt	r [RBP + l	ocal
11c1	MOV	RSI RAX		-
11c4	LEA	RDI,[s_The_he	eap_is_nea	r_0x
11cb	MOV	EAX,0X0		
11d0	CALL	<external>::</external>	orintf	
11d5	MOV	EAX,0X0		
11da	LEAVE			
11db	RET			

The Basics of Reverse Engineering What does Decompiled C look like?

- Denser than Assembly
- Some reasons it is difficult to read:
 - Missing variables names.
- Potential inaccuracies:
 - Missing arguments.

C _f	Decompile
1	
2	undefined
3	
4	{
5	undefin
6	void *l
7	
8	local_1
9	local_1
10	printf(
11	printf(
12	printf(
13	return
14	}
15	

: main – (pointers)

8 main(undefined4 param_1)

```
ed4 local_1c [3];
ocal_10;
```

```
.c[0] = param_1;
.0 = malloc(0x200);
"Main is at 0x%08llx.\n",main);
"The call stack is near 0x%08llx.\n",local_1c);
"The heap is near 0x%08llx.\n",local_10);
0;
```



A Quick Intro to Assembly Languages

- There are many of these languages, and they are different.
- This course focuses on x86 and will dabble in ARM.
 - ARM has three major dialects: ARM32, Thumb2, and ARM64
 - x86 has two major dialects: x86 and x86_64/amd64.
- You will be writing a little assembly, but reading a lot of it.

A Quick Intro to Assembly Languages JULIA EVANS Hexadecimal hexadecimal @bork

	nexade bas	ecimal se 16	is
<mark>0</mark> ~0	<mark>4</mark> →4	<mark>8</mark> ->8	с->1
1 ~>1	<mark>5~>5</mark>	<mark>9-></mark> 9	d ->1
<mark>2</mark> →2	<mark>6</mark> →6	<mark>a</mark> ->10	<mark>e</mark> -→1
<mark>3</mark> ∼ 3	<mark>7</mark> →7	b->11	f -→1
	bas	-16 b	ase 10

a 2-digit hex number is between 0 and 255 one byte 0x0 = 0 $0 \times 10 = 16$ is between 0x23 = 350 and 255 0xff = 255



powers of 2 are easier to recognize in hex

2²⁰ in decimal: 1048576 is that a power of 2? who knows! 2²⁰ in hex: 0x100000 more obviously a power of 2

case doesn't matter



things hexadecimal is used for

- → color codes! (eg 0xFF00FF)
- → memory addresses!
- displaying binary data! (like with hexdump)



A Quick Intro to Assembly Languages Registers

- Registers are like small variables that exist in hardware
- On x86_64 there are a lot of them but here are the most common
 - RAX, RBX, RCX, RDX, RSI, RDI, RBP, RSP, R8-R15, and RIP
 - These are all 64 bits in length.
 - commonly referred to as the Program Counter (PC)
- Floating point registers (3.14169)
- Flag Register (Zero, Signed, Carry, etc)
- Segment Registers (for memory stuff)
- AVX/SSE xmm/ymm/zmm 128/256/512 bits

See suggested reading list, Item [1]

• RIP (Instruction Pointer) is a special register that points to the next instruction to be executed. It is also



A Quick Intro to Assembly Languages Registers

- You can access subsets of the bits for many of them:
- RAX (all 64 bits) -> EAX (lower 32 bits) -> AX (lower 16 bits) -> AH (upper 8 bits of AX) -> AL (lower 8 bits of AX)

rax



A Quick Intro to Assembly Languages ADD RAX, RCX

- Each line is one Instruction.
- Each line begins with an Operation. In English grammar, a Verb.
- Parameters are typically called operands
 - 48 01 c8 is the machine code. ADD is often called the instruction mnemonic
- The first parameter is the Destination, storing the result.
- Instructions are grouped into Functions.
 - A function begins with the parameters on the Stack or in Registers.
 - A function ends with a standard instruction. (RET or BX LR.)

A Quick Intro to Assembly Languages ADD RAX, RCX

- This instruction is x86_64:
 - This specific syntax view of the instruction is called (Intel). There are other ways to represent it, such as Gas/AT&T syntax used by GNU tools (GCC)
 - Destination register comes first (with AT&T, it's the opposite! Why, oh why?)
 - The second register is one of the inputs.
 - The operation is ADDition.
- So what does this do?

A Quick Intro to Assembly Languages Common Operations

- Operations are unique to each assembly language, but some are common. • MOV, ADD, SUB, MUL -- Arithmetic

 - CALL, BL -- Function Calls
 - RET, BLR -- Function Returns
 - PUSH, POP -- Grow or shrink the Stack.
- A table can be handy for each new assembly language.
- Learn the common instructions, look up the rest.

A Quick Intro to Assembly Languages Stack

- You are probably familiar with the stack data type: Last in First Out (LIFO) • As opposed to Queue: First in First out (FIFO)
- The stack "grows down" from higher addresses to lower addresses
- Used to store local variables that were "statically allocated" at compile time
 - We say statically allocated because the size doesn't change when the program runs
- On x86_64 the stack is pointed to by RSP. It is an implicit operand in many instructions.

A Quick Intro to Assembly Languages PUSH

- Used to store data on the stack
- PUSH RAX
- Effectively
 - SUB RSP, 8
 - MOV [RSP], RAX
 - (* Note that "[]" denotes a dereference. Like var[8] = ## in C *)

A Quick Intro to Assembly Languages POP

- Used to take data from the stack and store it
- POP RAX
- Effectively
 - MOV RAX, [RSP]
 - ADD RSP, 8

- RSP := 0xFFF0
- RAX := 0xdeadbeef
- RCX := 0xd00dd00d

RSP 0xFFF0⁻

• PUSH RAX



• PUSH RCX



• POP RAX



• POP RCX



- RSP := 0x???
- RAX := 0x???
- RCX := 0x???

- RSP := 0xFFF0
- RAX := 0xd00dd00d
- RCX := 0xdeadbeef

A Quick Intro to Assembly Languages Control Flow

#include <stdio.h>

int main(int argc, char **argv){
 if(argc>5)
 printf("That's too many!\n");
 else
 printf("That's alright.\n");
}

A little C becomes a lot of Assembly. Sometimes it helps to think in a Graph.

A Quick Intro to Assembly Languages

Function Graph - main - 4 vertices (control)

Think of Control Flow as a Graph

	101
ur	nde
	.11
	.11
	.11
	.11
	.11
	.11
	.11
	.11

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A Quick Intro to Assembly Languages CMP RAX, RCX

- You can think of CMP as a signed subtraction
 - Signed uses the most significant bit (MSB) to indicate the sign. 1 == negative; 0 == positive
 - Uses twos complement
 - 8-bit char: 0xFF = -1; 0x01 = 1
- Effectively
 - temp = RAX RCX
- The value of temp is used to set the fields of the flag register. For example
- If temp == 0 then ZF = 1 (TRUE) else ZF = 0 (FALSE)
- If temp < 0 then SF = 1 else SF = 0
- There are other flags that may or may not be set depending on the instruction

A Quick Intro to Assembly Languages JMP & JXX

- JMP is an unconditional branching statement. It jumps where you tell it without checking any of the condition flags.
- JXX This is a family of instructions that jump if something is true. For example,
 - JZ 0x1234 will jump to 0x1234 if the Zero Flag is set to 1. This would happen if the operands of the last comparison were equal.
 - JNZ 0x1234 is the opposite, it will jump if not zero (ZF is set to 0)
- Different jump instructions check different flags but there are some that are equivalent.
 - JE is the same as JZ. If two operands are equal i.e. RAX and RCX then RAX RCX == 0.
- Lots of different types
 - JNE, JE, JG, JGE, JL, JLE, etc

A Quick Intro to Assembly Languages JMP & JXX

- Mentioned already but jumps are branching statements. i.e. they cause control flow to be non-linear
- For conditional branches there is a TRUE and a FALSE branch.
- Technically the FALSE branch just "falls through" (executes the next instruction
- The TRUE branch is taken if the check evaluates to TRUE
 - i.e. JZ evaluates the expression (ZF == 0). If ZF is equal to 0 then (ZF == 0)
 == 1

A Quick Intro to Assembly Languages JXX example. Branch taken?

mov	rax	02
mov	rcx	03
cmp	rax,	rc
jne	_end	



x5 x10 CX

A Quick Intro to Assembly Languages JXX example. Branch taken?

MOV	rax	0×
mov	rcx	0×
cmp	rax,	rc
je	_end	



x5 x10 CX

A Quick Intro to Assembly Languages JXX example. Branch taken?

MOV	rax	0 x
mov	rcx	0 x
cmp	rax,	rc
jl	_end	



x5 x10 CX

A Quick Intro to Assembly Languages CALL & RET

- What about functions?
- Just like in C programming we want to organise code so that it can be reused
- How to we get there and back again though?
- CALL 0x1234 (Typically a 5 byte instruction on x86_64)
- Effectively:
 - PUSH RIP + 5
 - JMP 0x1234

A Quick Intro to Assembly Languages CALL & RET

- Functions often take arguments
- x86 used the stack but x86_64 uses registers (mostly)
- Different calling conventions use different registers but for now we will focus on System V AMD64 ABI (used by the Linux-based Operating Systems)
 - RDI, RSI, RDX, RCX, R8, R10 (more than 6 uses the stack)
 - See: <u>https://en.wikipedia.org/wiki/X86_calling_conventions</u>
- The return value is stored in RAX

See suggested reading list, Item [2] re ABIs



A Quick Intro to Assembly Languages CALL & RET

- How do we get back
- RET (Return)
- Effectively:
 - POP RIP

A Quick Intro to Assembly Languages CALL / RET Example

- RIP := 0x1234
- RSP := 0xFFF0



A Quick Intro to Assembly Languages CALL / RET Example

- CALL 0x4320
 - RSP := 0xFFD8
 - RIP := 0x4320



A Quick Intro to Assembly Languages CALL / RET Example

- RET
 - RSP := 0xFFF0
 - RIP := 0x1239



A Quick Intro to Assembly Languages Final

- This was a firehose of information
- This isn't an assembly programming course
- There are a lot of instructions that you will have to look up on your own
- The INTEL instruction manual is your friend: 2A-2D
 - <u>https://www.intel.com/content/www/us/en/developer/articles/technical/</u> <u>intel-sdm.html</u>

Reversing a Simple Function

- Let's work out a few examples.
- First we'll see the assembly, and then we'll work backward to C.

Do this homework!

Reversing a Simple Function An Example

imul	rdi,	rc
mov	rax,	rc
ret		



di di

Reversing a Simple Function An Example

- imul rdi, rdi rax, rdi MOV ret
- Register rdi is the first parameter
- "imul rdi, rdi" MULtiplies RDI by itself and stores the result in RDI.
- "mov rax, rdi" MOVes the value of RDI into rax.
- "ret" is the standard return function on x86_64.

Reversing a Simple Function An Example

imul rdi, rdi mov rax, rdi ret

int square(int num) { return num * num;



Reversing a Simple Function Another Example

0	x0	0	0	0
0	x0	0	0	3
0	x0	0	0	6
0	x0	0	0	8
0	x0	0	0	b

mov cmp jle mov ret

rax,	rdi
rax,	rsi
$0 \times 0 0 0$	b
rax,	rsi

Reversing a Simple Function Another Example

$0 \times 0 0 0 \times 0$	mov	rax, rdi
0x0003	cmp	rax, rsi
0x0006	jle	0x000b
8000x0	mov	rax, rsi
0x000b	ret	

- Register rdi is the first parameter, rsi is the second
- "mov, rax, rdi" sets rax equal to rdi. This is an optimisation to save a branch.
- "cmp rax, rsi" Compares rax (which was rdi) to rsi
- "jle" Jumps if rax is less than or equal to rsi
- "mov, rax, rsi" sets rax equal to rsi.
- "ret" is the standard return function on x86_64.

Reversing a Simple Function Another Example

 $0 \times 0 0 0 \times 0$ 0x0003 0x0006 8000x00x000b

MOV mov ret

rax, rdi cmp rax, rsi jle 0x000b rax, rsi

int min(int a, int b) { if (a <= b) { return a; else return b;

Useful Tools for This Course

- GNU Objdump -- Command-line disassembler for many architectures.
- <u>ghidra-sre.org</u> -- GHIDRA, the NSA's reverse engineering tool.
- radare.org -- A free, command-line reverse engineering toolkit.
- <u>godbolt.org</u> -- Compiler Explorer, a tool to view assembly from snippets of C.
- GDB, LLDB -- Debuggers are very handy for exploring samples.
- Pen and Paper! -- Reverse engineering is puzzle solving.



Disassembling a Binary with Objdump

dell% objdump –d first

first: file format elf64-x86-64

Disassembly of section .init:

0000000000001000 <_init>: 1000:48 83 ec 08 \$0x8,%rsp sub 1004:48 8b 05 dd 2f 00 00 mov 0x2fdd(%rip),%rax # 3fe8 <__gmon_start__@Base> 100b:48 85 c0 %rax,%rax test 1012 <_init+0x12> 100e:74 02 je 1010:ff d0 call *%rax 1012:48 83 c4 08 add \$0x8,%rsp 1016: c3 ret

Disassembling with GDB

dell% gdb first Reading symbols from first... (No debugging symbols found in first) (gdb) disassemble __init Dump of assembler code for function __init: 0x0000000000001000 <+0>:sub \$0x8,%rsp 0x2fdd(%rip),%rax 0x0000000000000004 <+4>:mov 0x000000000000100b <+11>:test %rax,%rax 0x000000000000100e <+14>:je **0**x1012 < init+18> 0x0000000000001010 <+16>:call *%rax 0x0000000000001012 <+18>:add \$0x8,%rsp 0x0000000000001016 <+22>:ret End of assembler dump. (gdb)



0x3fe8

Disassembling with Radare2

dell% r2 first [0x00001050] > pd 7 @0x1000 ;-- section..init: ;-- segment.LOAD1: ;-- _init: 0x00001000 4883ec08 sub rsp, 8 0x00001004 488b05dd2f00 0x0000100b 4885c0 -> 0x00001012 4883c408 **0x00001016** c3 [0x00001050]>

```
mov rax, qword [reloc.__gmon_start] ; [0x3fe8:8]=0
test rax, rax
je 0x1012
call rax
add rsp, 8
ret
```



Decompiling with GHIDRA

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<u>File Edit Analysis Graph Navigation Search Select Tools</u>	s <u>W</u> indow <u>H</u> elp							
🔚 (
Program Trees 🔂 🔂 🔭 🗙	🖽 Listing: first				r 🗈 🔁 🖳 🛱 🖌	💩 📑 - 🗙	€ Decompile: _init - ((first) 🧐 🔓 🛛
<pre>inii itext i.text i.plt.got init i.init i.rela.plt i.rela.dyn i.rela.dyn i.gnu.version_r i.gnu.version i.dynstr i.gnu.hash i.onte.ABI-tag i.note.gnu.build-id intere Program Tree × fdo_global_dtors_aux fgmon_start flibc_csu_fini flibc_csu_init flibc_csu_init</pre>		<pre>// ram:001010 // ******************************</pre>	00-ram:00101016 ************************************	<pre>************************************</pre>	<pre>*** * Entry Point(*), 001000f8(*),libc_csu_init:00101188(c), 00103e10(*), _elfSectionHeaders::000002d0(*) *t] = 00105018 = ?? undefinedgmon_start() 0010100e(j) *** ***</pre>		<pre>1 2 int _init(EVP_P+ 3 4 5 int iVar1; 6 7 8 9 10 10 10 10 10 10 10 10 10 10 10 10 10</pre>	<pre>EY_CTX *ctx) _start();</pre>
start								
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						0010100b	_init	TEST RAX,RAX





Day 1 Homework

- From disassembly to pseudocode:
 - Take the listing you are given and write the C-like pseudocode
 - What will the function return given a specific value?
- For next class, have Ghidra installed and make sure that you can ssh into the Babylon servers from your laptop
 - <u>https://kb.thayer.dartmouth.edu/article/361-linux-services</u>
 - For login with Kerberos tokens rather than endlessly retyping your password: <u>https://hackmd.io/e5Ft3DXCRze6NGnOudCt4Q</u>
 - Also, enable GSSAPI in your .ssh/config for passwordless SCP to work: https://services.dartmouth.edu/TDClient/1806/Portal/KB/ArticleDet?ID=89203